

Express Mail Mailing Label No. EL750476114US

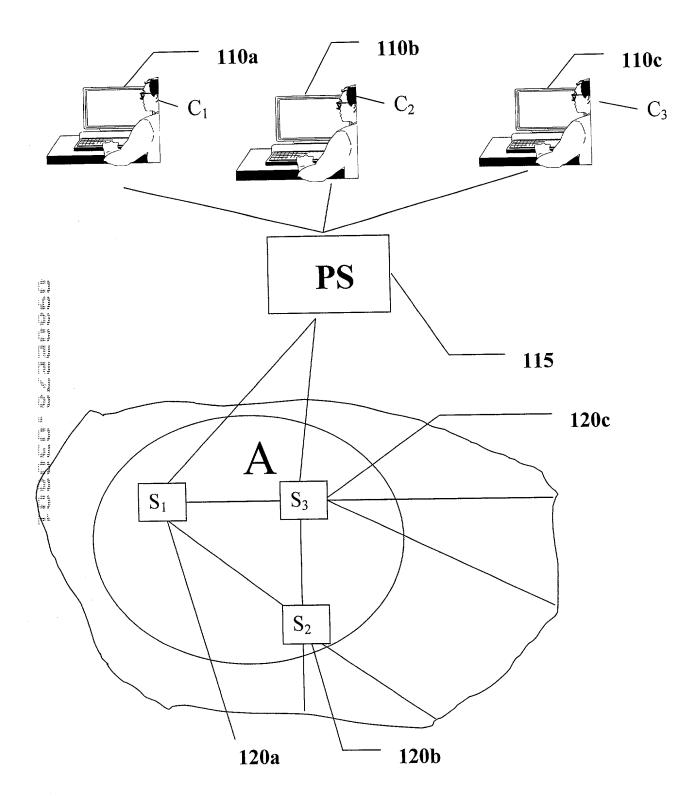


FIGURE 2

Title: Targeted Delivery of Informational Content with Privacy Protection Inventor: Juels Serial No. Not yet assigned Atty Docket No. RSA-044 (7216/66) Atty/Agent: Ira V. Heffan Express Mail Mailing Label No. EL750476114US

$$C_1, C_2, C_3 \leftarrow A$$
 STEP 31

$$C_i$$
 computes  $r_i = f(P_{C_i})$  STEP 32

$$C_i \xrightarrow{r_i} A$$
 STEP 33

$$C_i \leftarrow A$$
 STEP 35

$$C_1, C_2, C_3 \stackrel{f}{\longleftarrow} PS \stackrel{f}{\longleftarrow} A$$
 STEP 41

$$C_1$$
 computes  $r_1 = f(P_{C1})$   
 $C_2$  computes  $r_2 = f(P_{C2})$   
 $C_3$  computes  $r_3 = f(P_{C3})$ 

$$C_{1} \xrightarrow{r_{1}} PS$$

$$C_{2} \xrightarrow{r_{2}} PS$$

$$C_{3} \xrightarrow{r_{3}} PS$$

$$STEP 43$$

$$(x_1, r_1) (x_2, r_2) (x_3, r_3)$$
PS  $\longrightarrow$  A STEP 44

r<sub>i</sub> causes A to select ad<sub>ri</sub>

$$(x_1, ad_1) (x_2, ad_2) (x_3, ad_3)$$
PS  $\longrightarrow$  A STEP 46

**STEP 45** 

$$C_1, C_2, C_3 \stackrel{ad_i}{\longleftarrow} PS$$
 STEP 47

## FIGURE 4

f

STEP 51 
$$C_1, C_2, C_3 \leftarrow A$$

STEP 52 
$$C_1$$
 computes  $r_1 = f(P_{C1})$  and encrypts  $E_y[r_1]$   $C_2$  computes  $r_2 = f(P_{C2})$  and encrypts  $E_y[r_2]$   $C_3$  computes  $r_3 = f(P_{C3})$  and encrypts  $E_y[r_3]$ 

STEP 53
$$C_{1} \xrightarrow{\{E_{y}[r_{1}], x_{1}\}} BB$$

$$C_{2} \xrightarrow{\{E_{y}[r_{2}], x_{2}\}} BB$$

$$C_{3} \xrightarrow{\{E_{y}[r_{3}], x_{3}\}} BB$$

STEP 54 Servers collect 
$$V_1 = \{ E_y[r_i], x_i \}_{i=1}^k$$

STEP 55 Servers mix 
$$V_1$$
 by random secret permutation  $\sigma_1$  to obtain  $V_2 = \{r_{\sigma I}(i), E_y[\sigma_1(i)]\}_{i=1}^k$ 

STEP 56 Servers replace each 
$$r_j$$
 in  $V_2$  with  $ad_{rj}$  to obtain  $V'_2 = \{ad_r, E_y[\sigma_1(i)]\}_{i=1}^k$ 

STEP 57 Servers mix V'<sub>2</sub> by random secret permutation 
$$\sigma_2$$
 to obtain V<sub>3</sub> = { $(E_y[ad_{\sigma_2(i)}], \sigma_2(i)$ } $_{i=1}^k$ 

STEP 58 Servers apply quorum controlled asymmetric proxy re-encryption to obtain 
$$V_4 = (E_{yci}[ad_{ri}], i)_{i=1}^k$$

STEP 59 
$$C_1, C_2, C_3 \leftarrow A$$

STEP 60 
$$C_1$$
,  $C_2$ ,  $C_3$  decrypt  $E_{yci}[ad_{ri}]$  to receive  $ad_{ri}$ 

## FIGURE 5

f

		<b>A</b>
STEP 61	C -	$\mathbf{A}$

STEP 62 C computes 
$$r = f(P)$$
 and encrypts  $E_y[r]$ 

STEP 63 
$$C \xrightarrow{E_y[r_1]} BB$$

**STEP 64** Servers encrypt 
$$ad_i$$
 to generate  $U_1 = \{(j, E_y[ad_j])\}_{j=1}^n$ 

STEP 65 Servers mix 
$$U_1$$
 by random secret permutation  $\sigma$  to obtain  $U_2 = (E_y[\sigma(j)], E_y[ad_{\sigma(j)}])^n_{j=1}$ 

**STEP 66** Servers perform a distributed plaintext equality test to find 
$$E_{\nu}[j] \sim E_{\nu}[r]$$
 and obtain  $U_3 = (E_{\nu}[r], E_{\nu}[ad_r])$ 

STEP 68 
$$C \leftarrow A$$

## FIGURE 6